1 Introduction to Mathematics (Continued)

1.6 And Even More on Quantification

Before we get started, it is important to remind you that we will not be explicitly using the symbolic representation of a given statement in terms of quantifiers and logical connectives. Nonetheless, having this notation at our disposal allows us to compartmentalize the abstract nature of mathematical propositions and will provide us with a way to talk about the meta-concepts surrounding the construction of proofs.

Definition 1.80. Two quantified propositions are **equivalent in a given universe of discourse** iff they have the same truth value in that universe. Two quantified propositions are **equivalent** iff they are equivalent in every universe of discourse.

Exercise 1.81. Consider the propositions $(\forall x)(x > 3)$ and $(\forall x)(x \ge 4)$. Are these propositions equivalent if the universe of discourse is the set of integers? (*Hint:* What are their truth values in this case?) Come up with two different universes of discourse that yield different truth values for these propositions. What can you conclude?

Remark 1.82. At this point it is worth pointing out an important distinction. Consider the propositions "All cars are red" and "All natural numbers are positive". Both of these are instances of the **logical form** $(\forall x)P(x)$. It turns out that the first proposition is false and the second is true; however, the logical form is neither true or false. A logical form is a blueprint for particular propositions. If we are careful, it makes sense to talk about whether two logical forms are equivalent. For example, $(\forall x)(P(x) \implies Q(x))$ is equivalent to $(\forall x)(\neg Q(x) \implies \neg P(x))$. For fixed P(x) and Q(x), these two forms will always have the same truth value independent of the universe of discourse. If you change P(x) and Q(x), then the truth value may change, but the two forms will still agree.

Theorem 1.83. Let P(x) be a predicate. Then

- 1. $\neg(\forall x)P(x)$ is equivalent to $(\exists x)(\neg P(x));$
- 2. $\neg(\exists x)P(x)$ is equivalent to $(\forall x)(\neg P(x))$.

Exercise 1.84. Negate each of the following. Disregard the truth value and the universe of discourse.

- 1. $(\forall x)(x > 3)$
- 2. $(\exists x)(x \text{ is prime } \land x \text{ is even})$
- 3. All cars are red.
- 4. Every wookie is named Chewbacca.
- 5. Some hippies are republican.
- 6. For all $x \in \mathbb{N}$, $x^2 + x + 41$ is prime.
- 7. There exists $x \in \mathbb{Z}$ such that $1/x \notin \mathbb{Z}$.
- 8. There does not exist a function f such that if f is continuous, then f is not differentiable.

Using Theorem 1.83 and our previous results involving quantification, we can negate complex mathematical propositions by working from left to right.

Example 1.85. Consider the proposition

$$(\exists x \in \mathbb{R}) (\forall y \in \mathbb{R}) (x + y = 0).$$

It turns out that this statement is false, which means that its negation is true. That is,

$$\neg (\exists x \in \mathbb{R}) (\forall y \in \mathbb{R}) (x + y = 0),$$

which is equivalent to

$$(\forall x \in \mathbb{R}) (\exists y \in \mathbb{R}) (x + y \neq 0)$$

is true.

Example 1.86. Consider the proposition

$$(\forall x)[x > 0 \implies (\exists y)(y < 0 \land xy > 0)]$$

(which happens to be false). Then

$$\neg(\forall x)[x > 0 \implies (\exists y)(y < 0 \land xy > 0)]$$

is equivalent to

 $(\exists x)[x > 0 \land \neg (\exists y)(y < 0) \land xy > 0)],$

which is equivalent to

$$(\exists x)[x > 0 \land (\forall y)(y \ge 0 \lor xy \le 0)].$$

Exercise 1.87. What previous theorems were used when negating the proposition in the previous example.

Exercise 1.88. Negate each of the following. Disregard the truth value and the universe of discourse.

- 1. $(\forall n \in \mathbb{N}) (\exists m \in \mathbb{N}) (m < n)$
- 2. $(\forall x, y, z \in \mathbb{Z})((xy \text{ is even} \land yz \text{ is even}) \implies xy \text{ is even})$
- 3. For all positive real numbers x there exists a real number y such that $y^2 = x$.
- 4. There exists a married person x such that for all married people y, x is married to y.

At this point, we should be able to use our understanding of quantification to construct counterexamples to complicated false propositions and proofs of complicated true propositions. Here are some general proof structures for various logical forms.

Direct Proof of $(\forall x)P(x)$:

Proof. Let x be arbitrary object in universe of discourse.

(Use definitions and previous results.)

Therefore, P(x) is true. Since x was arbitrary, for all x, P(x).

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This work is an adaptation of notes written by Stan Yoshinobu of Cal Poly and Matthew Jones of California State University, Dominguez Hills.

Proof of $(\forall x)P(x)$ by Contradiction:

Proof. For sake of a contradiction, assume that there exists x in the universe of discourse such that $\neg P(x)$.

(Do something to derive a contradiction.)

This is a contradiction. Therefore, for all x, P(x) is true.

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Direct Proof of $(\exists x)P(x)$:

Proof. (Either use definitions and previous results to deduce that an x exists such that P(x) is true or if you think you have an x that works, just verify that it does.)

(Do stuff.)

Therefore, there exists x such that P(x).

Proof of $(\exists x)P(x)$ by Contradiction:

Proof. For sake of a contradiction, assume that for all x, $\neg P(x)$.

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(Do something to derive a contradiction.)

This is a contradiction. Therefore, there exists x such that P(x).

Question 1.89. Suppose P(x) is a proposition such that $(\forall x)P(x)$ is false. Which of the above proof situations is identical to providing a counterexample to this proposition?

Remark 1.90. It is important to point out that sometimes we will have to combine various proof techniques in a single proof. For example, if you wanted to prove a proposition of the form $(\forall x)(P(x) \implies Q(x))$ by contradiction, we would start by assuming that there exists x such that P(x) and $\neg Q(x)$.

Problem 1.91. For each of the following statements, determine its truth value. If the statement is false, provide a counterexample. Prove at least two of the true statements.

- 1. For all $n \in \mathbb{N}$, $n^2 \ge 5$.
- 2. There exists $n \in \mathbb{N}$ such that $n^2 1 = 0$.
- 3. There exists $x \in \mathbb{N}$ such that for all $y \in \mathbb{N}$, $y \leq x$.
- 4. For all $x \in \mathbb{Z}, x^3 \ge x$.
- 5. For all $n \in \mathbb{Z}$, there exists $m \in \mathbb{Z}$ such that n + m = 0.
- 6. There exists integers a and b such that 2a + 7b = 1.
- 7. There do not exist integers m and n such that 2m + 4n = 7.
- 8. For all integers a, b, c, if a divides bc, then either a divides b or a divides c.

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