

Braid graphs in simply-laced triangle-free Coxeter systems are partial cubes

Part 1

ACGT

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November 2021

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Definition

A **Coxeter system** consists of a group W (called a **Coxeter group**) generated by a set S of involutions with presentation

$$W = \langle S \mid s^2 = e, \quad (st)^{m(s,t)} = e \rangle,$$

where $m(s, t) \geq 2$ for $s \neq t$.

Comments

- The elements of S are distinct as group elements.
- $m(s, t)$ is the order of st .

Coxeter Systems

Since s and t are involutions, the relation $(st)^{m(s,t)} = e$ can be rewritten:

$$m(s, t) = 2 \implies st = ts \quad \left. \vphantom{m(s, t) = 2} \right\} \text{commutation relations}$$

$$\left. \begin{array}{l} m(s, t) = 3 \implies sts = tst \\ m(s, t) = 4 \implies stst = tsts \\ \vdots \end{array} \right\} \text{braid relations}$$

This allows the replacement

$$\underbrace{sts \cdots}_{m(s,t)} \mapsto \underbrace{tst \cdots}_{m(s,t)}$$

in any word, which is called a **commutation move** if $m(s, t) = 2$ and a **braid move** if $m(s, t) \geq 3$.

Definition

We can encode (W, S) with a unique Coxeter graph Γ having:

- Vertex set = S
- $\{s, t\}$ edge labeled with $m(s, t)$ whenever $m(s, t) \geq 3$

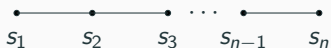
Comments

- Typically labels of $m(s, t) = 3$ are omitted.
- Edges correspond to non-commuting pairs of generators.
- Given Γ , we can uniquely reconstruct the corresponding (W, S) .

Coxeter Systems of Type A

Example

The Coxeter system of type A_n is defined by the following graph.



Then $W(A_n)$ is subject to:

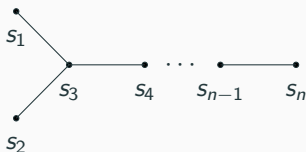
- $s_i^2 = e$ for all i
- $s_i s_j = s_j s_i$ if $|i - j| > 1$
- $s_i s_j s_i = s_j s_i s_j$ if $|i - j| = 1$.

In this case, $W(A_n)$ is isomorphic to the symmetric group S_{n+1} under the correspondence $s_i \mapsto (i, i + 1)$.

Coxeter Systems of Type D

Example

The Coxeter system of type D_n is defined by the following graph.



Then $W(D_n)$ is subject to:

- $s_i^2 = e$ for all i
- $s_i s_j = s_j s_i$ when $|i - j| > 1$ and $3 \notin \{i, j\}$
- $s_i s_3 s_i = s_3 s_i s_3$ for $i \in \{1, 2, 4\}$
- $s_i s_j s_i = s_j s_i s_j$ when $|i - j| = 1$ and $i, j \in \{4, 5, \dots, n\}$.

The group $W(D_n)$ is isomorphic to the index 2 subgroup of the group of signed permutations on n letters having an even number of sign changes.

Reduced Expressions & Matsumoto's Theorem

Definition

A word $\alpha = s_{x_1} s_{x_2} \cdots s_{x_m} \in S^*$ is called an **expression** for w if it is equal to w when considered as a group element. If m is minimal among all expressions for w , α is called a **reduced expression**.

$\mathcal{R}(w)$ = set of reduced expressions for w

Reduced Expressions & Matsumoto's Theorem

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Matsumoto's Theorem

Any two reduced expressions for $w \in W$ differ by a sequence of commutation & braid moves.

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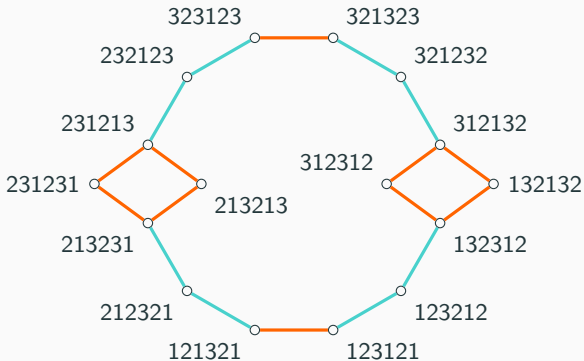
For $w \in W$, define the **Matsumoto graph** $\mathcal{M}(w)$ via:

- Vertex set = $\mathcal{R}(w)$
- $\{\alpha, \beta\}$ edge iff α and β are related via a single **commutation** or **braid** move

Matsumoto Graph

Example

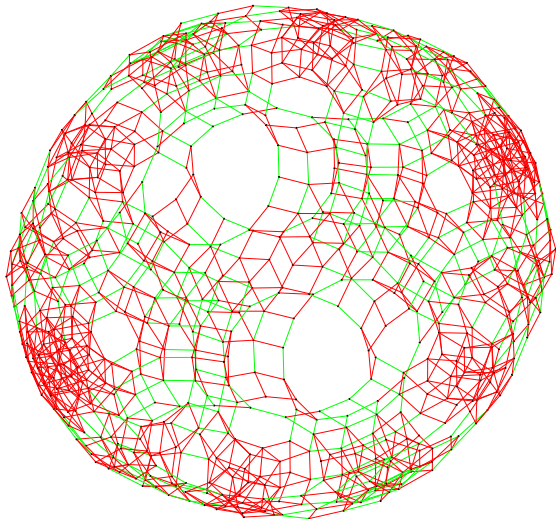
Consider the reduced expression $\alpha = 121321$ for $w \in W(A_3)$. Then $\mathcal{M}(w)$ is as follows:



Matsumoto Graph

Example

Here is the Matsumoto graph for the longest element in type A_4 .



Braid Equivalence & Braid Graphs

Definition

If $\alpha, \beta \in \mathcal{R}(w)$, then α and β are **braid equivalent** iff α and β are related by a sequence of braid moves.

Comments

- Braid equivalence is an equivalence relation.
- Equivalence classes are called **braid classes**, denoted $[\alpha]$.

Braid Equivalence & Braid Graphs

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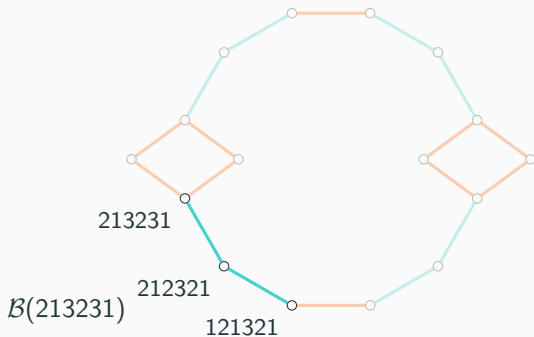
Definition

We can encode a braid class $[\alpha]$ in a **braid graph**, denoted $\mathcal{B}(\alpha)$:

- Vertex set = $[\alpha]$
- $\{\gamma, \beta\}$ edge iff γ and β are related via a single **braid move**

Braid graphs are the maximal **turquoise** connected components in the Matsumoto graph. Not to be confused with contracting the braid edges of a Matsumoto graph.

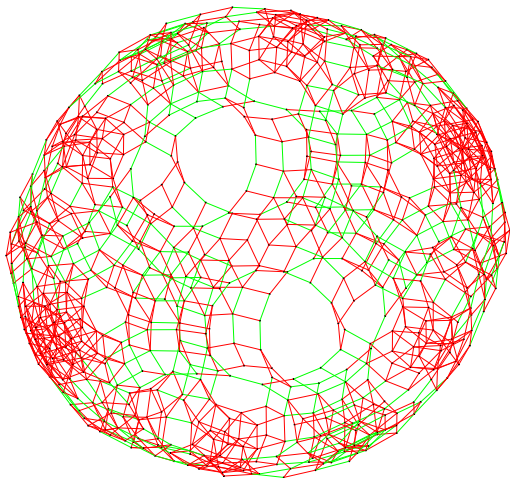
Example



Braid Graphs

Example

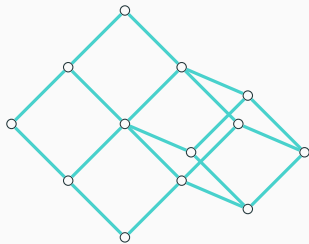
Each of the maximal **green** connected components in the following Matsumoto graph is a braid graph corresponding to a braid class.



Braid Graphs

Example

Consider the reduced expression $\alpha = 31323431323$ for some $w \in W(D_4)$. Then $\mathcal{B}(\alpha)$ is as follows, where α is the vertex of degree 5.



Definition

Suppose α is a reduced expression for $w \in W$ consisting of m letters. Loosely speaking, α is **link** if there is a sequence of overlapping braid moves that “cover” the positions $1, 2, \dots, m$. If α is a link, then the corresponding braid class $[\alpha]$ is called a **braid chain**.

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Example

Consider the reduced expression $\alpha = 343546576$ for some $w \in W(A_7)$.

343546576 **43454**6576 43**54565**76 4354**65676** 435465**767**



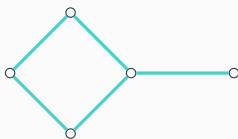
In this example, every reduced expression is a link and the braid class is a braid chain.

Example

Let $\alpha = 3134323$ be a reduced expression for some $w \in W(D_4)$.

$$\overline{3134323}$$

Then α is a link and $[\alpha]$ is a braid chain. The corresponding braid graph is as follows, where α is the vertex of degree 3.



Example

Now, let $\alpha = 1213243676$ be a reduced expression for some $w \in W(A_7)$. It turns out that α is not a link, but rather a product of two links.

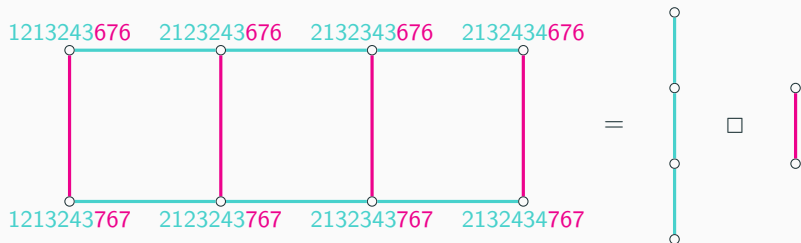
$$1213243 \mid 676$$

Links & Braid Chains

Example

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1213243 | 676



Comments

- Every reduced expression factors uniquely into maximal links, called a **braid link factorization**.
- Describing the maximal links and their corresponding braid chains is tricky business!
- We have a nice characterization for triangle-free simply-laced Coxeter systems.

Braid Graphs for Braid Factorizations

Theorem

If α is a reduced expression for $w \in W$ having braid link factorization

$$\alpha = \beta_1 | \beta_2 | \cdots | \beta_m,$$

then $\mathcal{B}(\alpha)$ is the box product of the braid graphs for each β_i .

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Comment

- The upshot is that if you want to understand the structure of braid graphs, you must first characterize the braid graphs for links.
- We've classified the braid graphs for links in types A_n and D_n (and others at least "in my head").
- In the case of type A_n , links have odd length and the corresponding braid graphs are paths.

Braid Graphs for Braid Factorizations

Theorem (Fisher et al. \rightarrow Bidari & Ernst)

If α is a reduced expression for $w \in W(A_n)$ having braid link factorization

$$\alpha = \beta_1 | \beta_2 | \cdots | \beta_m$$

such that each factor has $2k_i - 1$ generators, then

$$\mathcal{B}(\alpha) = \left. \begin{array}{c} \circ \\ | \\ \circ \\ \vdots \\ \circ \end{array} \right\} k_1 \square \left. \begin{array}{c} \circ \\ | \\ \circ \\ \vdots \\ \circ \end{array} \right\} k_2 \square \cdots \square \left. \begin{array}{c} \circ \\ | \\ \circ \\ \vdots \\ \circ \end{array} \right\} k_m$$

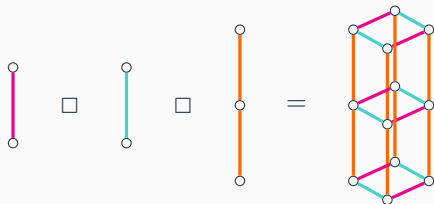
Braid Graphs for Braid Factorizations

Example

Consider the following braid link factorization for a reduced expression for an element in $W(A_7)$.

$$\alpha = 121 \mid 434 \mid 65676$$

The resulting braid graph is shown below:



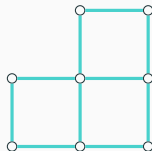
Fibonacci Links in Type D

Consider the Coxeter system of type D_4 . Let $\{a, b, c\} = \{1, 2, 4\}$. Every reduced expression that is braid equivalent to one of the following is called a **Fibonacci link** (in type D_4). The corresponding braid graph is depicted on the right.

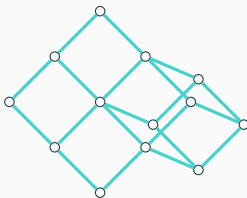
$3a3b3c3$



$3a3b3c3a3$



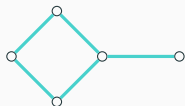
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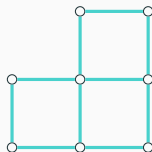
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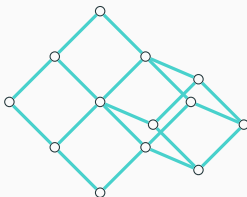
$3a3b3c3$



$3a3b3c3a3$



$3c3a3b3c3a3$



Each one of the graphs above corresponds to a **Fibonacci cube graph**!

Classification of Braid Graphs for Links in Type D

Theorem

In type D_n , every link is braid equivalent to either a “type A ” link or a “type A extension” of a Fibonacci link. As a consequence, braid graphs for links in type D_n are either paths or “type A extensions” of braid graphs for Fibonacci links.

Choices for a, b, c determine whether we can extend; need 343 on an end.

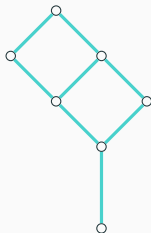
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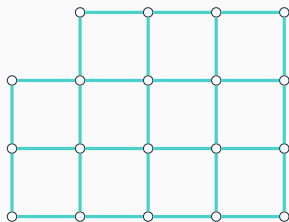
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Examples



453431323



453431323435465

Classification of Braid Graphs in Type D

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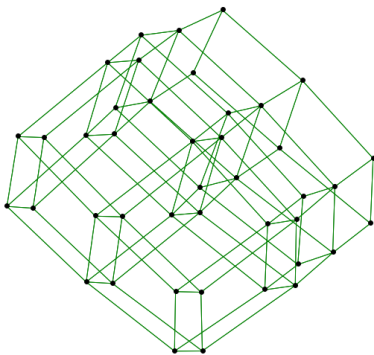
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Example



$$\alpha = 453431323 \mid 56576 \mid 898$$