

# Classification of the T-avoiding permutations

J. Cormier, Z. Goldenberg, J. Kelly, C. Malbon  
Directed by D.C. Ernst

Plymouth State University  
Mathematics Department

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# The symmetric group

## Definition

The **symmetric group**  $S_n$  is the collection of bijections from  $\{1, 2, \dots, n\}$  to  $\{1, 2, \dots, n\}$  where the operation is function composition (left  $\leftarrow$  right). Each element of  $S_n$  is called a **permutation**.

## Comment

We can think of  $S_n$  as the group that acts by rearranging  $n$  coins.

One way of representing permutations is via **cycle notation**, which we will illustrate by way of example.

## Example

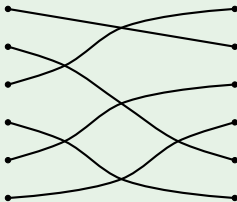
Consider  $\sigma = (1\ 3\ 5\ 2)(4\ 6)$ . This means  $\sigma(1) = 3$ ,  $\sigma(3) = 5$ ,  $\sigma(5) = 2$ ,  $\sigma(2) = 1$ ,  $\sigma(4) = 6$ , and  $\sigma(6) = 4$ .

# String diagrams

A second way of representing permutations is via **string diagrams**, which we again introduce by way of example.

## Example

Consider  $\sigma = (1\ 3\ 5\ 2)(4\ 6)$  from previous example.



## Comment

Given a permutation  $\sigma$ , there are many ways to draw the associated string diagram. However, we adopt the following conventions:

1. no more than two strings cross each other at a given point,
2. strings are drawn so as to minimize crossings.

$S_n$  is generated by the adjacent 2-cycles:

$$(1\ 2), (2\ 3), \dots, (n-1\ n).$$

That is, every element of  $S_n$  can be written as a product of the adjacent 2-cycles.

Define

$$s_i = (i\ i+1)$$

so that  $s_1, s_2, \dots, s_{n-1}$  generate  $S_n$ .

## Comments

$S_n$  satisfies the following relations:

1.  $s_i^2 = 1$  for all  $i$  (2-cycles have order 2)
2. short braid relations:  $s_i s_j = s_j s_i$ , for  $|i - j| \geq 2$  (disjoint cycles commute)
3. long braid relation:  $s_i s_j s_i = s_j s_i s_j$ , for  $|i - j| = 1$ .

## Definition

If  $s_{x_1} s_{x_2} \cdots s_{x_m}$  is an expression for  $\sigma \in S_n$  and  $m$  is minimal, then we say that the expression is **reduced**.

## Example

Consider  $\sigma = s_2 s_1 s_2 s_3 s_1 s_2 \in S_4$ . We see that

$$s_2 s_1 s_2 s_3 s_1 s_2 = s_1 s_2 s_1 s_3 s_1 s_2 = s_1 s_2 s_1 s_3 s_1 s_2 = s_1 s_2 s_1 s_1 s_3 s_2 = s_1 s_2 s_1 s_1 s_3 s_2 = s_1 s_2 s_3 s_2.$$

So, the original expression was *not* reduced, but it turns out that the last expression on the right is reduced.

## Theorem (Matsumoto)

*Any two reduced expressions for  $\sigma \in S_n$  differ by a sequence of braid relations.*

## Example (continued)

The only reduced expressions for  $\sigma$  are:  $s_1 s_2 s_3 s_2$ ,  $s_1 s_3 s_2 s_3$ , and  $s_3 s_1 s_2 s_3$ .

# Heaps

A third way of representing permutations is via **heaps**. Fix a reduced expression  $s_{x_1} s_{x_2} \cdots s_{x_m}$  for  $\sigma \in S_n$ . Loosely speaking, the heap for this expression is a set of lattice points (called **nodes**), one for each  $s_{x_i}$ , embedded in  $\mathbb{N} \times \mathbb{N}$  such that:

- The node corresponding to  $s_{x_i}$  has vertical component equal to  $n + 1 - x_i$  (smaller numbers at the top),
- If  $i < j$  and  $s_{x_i}$  does not commute with  $s_{x_j}$ , then  $s_{x_i}$  occurs to the left of  $s_{x_j}$ .

## Example

Consider  $s_1 s_2 s_3 s_2$ ,  $s_1 s_3 s_2 s_3$ , and  $s_3 s_1 s_2 s_3$  from the previous example. It turns out, there are two distinct heaps.

$s_1$  ●

$s_2$  ●

$s_2$  ●

$s_3$  ●

and

$s_1$  ●

$s_2$  ●

$s_3$  ●

$s_3$  ●

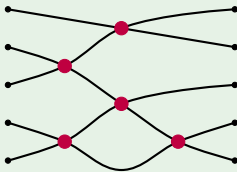
## Comment

If two reduced expressions for  $\sigma$  differ by a sequence of short braid relations, then they have the same heap. In particular, if no reduced expression contains an opportunity to provide a long braid relation, then  $\sigma$  has a unique heap.

The points at which two strings cross correspond to nodes in the heap. Hence, we may overlay strings on top of a heap by drawing the strings from right to left so that they cross at each entry in the heap where they meet and bounce at each lattice point not in the heap.

Conversely, each string diagram corresponds to a heap by taking all of the points where the strings cross as the nodes of the heap.

## Example

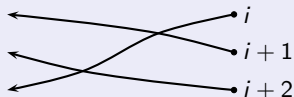


## Definition

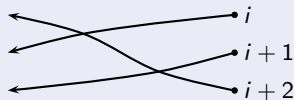
We say that a permutation  $\sigma$  has **Property T** iff there exists  $i$  such that either

1.  $\sigma(i) > \sigma(i+1), \sigma(i+2)$ ,

2.  $\sigma(i+2) < \sigma(i), \sigma(i+1)$ .

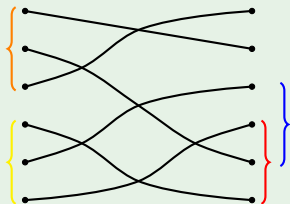


or



## Example

Consider the following permutation  $\sigma = (1\ 3\ 5\ 2)(4\ 6)$ .



We see that  $\sigma$  and  $\sigma^{-1}$  each have Property T in two spots.



## Definition

We say that a permutation  $\sigma$  is **T-avoiding** iff neither  $\sigma$  or  $\sigma^{-1}$  has Property T.

## Theorem (CEGKM)

*A permutation  $\sigma$  is T-avoiding iff  $\sigma$  is a product of disjoint adjacent 2-cycles.*

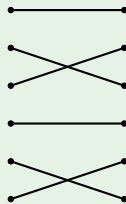
## Sketch of proof

Fix a reduced expression for  $\sigma$ , say  $s_{x_1} s_{x_2} \cdots s_{x_m}$ , and consider the heap for this reduced expression. The reverse implication of the theorem is trivial. For the forward direction, consider the contrapositive:

If  $\sigma$  is not a product of disjoint adjacent 2-cycles, then  $\sigma$  or  $\sigma^{-1}$  has Property T.

## Example

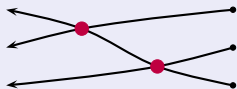
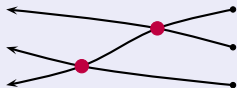
The permutation  $\sigma = (2\ 3)(5\ 6)$  is T-avoiding.



## The easy case

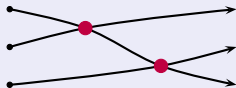
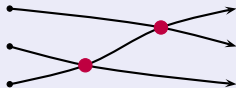
The easy case occurs when a node in the second column on either side is "blocked" by at most one node in the first column.

1.  $\sigma$  has the property:

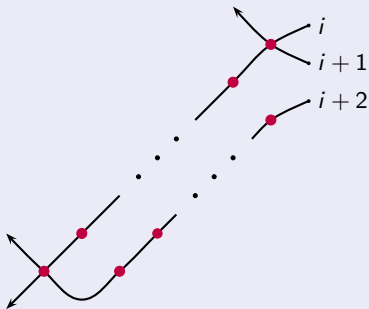


or

2.  $\sigma^{-1}$  has the property:



## The hard case



## Theorem

*By applying a sequence of long braid relations, you can convert a heap in the hard case to a heap in the easy case.*

The previous theorem provides the necessary motivation for generalizing the definition of Property T in other Coxeter groups.

## Definition

A **Coxeter group** consists of a group  $W$  together with a generating set  $S$  consisting of elements of order 2 with presentation

$$W = \langle S : s^2 = 1, (st)^{m(s,t)} = 1 \rangle,$$

where  $m(s, t) \geq 2$  for  $s \neq t$ .

## Comment

Since  $s$  and  $t$  are elements of order 2, the relation  $(st)^{m(s,t)} = 1$  can be rewritten as

$$m(s, t) = 2 \quad \Longrightarrow \quad st = ts \quad \left. \vphantom{m(s, t) = 2} \right\} \text{short braid relations}$$

$$m(s, t) = 3 \quad \Longrightarrow \quad sts = tst$$

$$m(s, t) = 4 \quad \Longrightarrow \quad stst = tsts$$

$$\vdots$$
$$\left. \vphantom{m(s, t) = 4} \right\} \text{long braid relations}$$

## Example (Type A)

The symmetric group  $S_{n+1}$  with the adjacent 2-cycles as a generating set is a Coxeter group of type  $A_n$ .

## Example (Type B)

Coxeter groups of type  $B_n$  ( $n \geq 2$ ) are defined by:

1.  $s_i^2 = 1$  for all  $i$ ,
2.  $s_i s_j = s_j s_i$  if  $|i - j| > 1$ ,
3.  $s_i s_j s_i = s_j s_i s_j$  if  $|i - j| = 1$  and  $1 < i, j \leq n$ ,
4.  $s_1 s_2 s_1 s_2 = s_2 s_1 s_2 s_1$ .

## Definition

Let  $(W, S)$  be a Coxeter group and let  $w \in W$ . Then  $w$  has **Property T** iff  $w$  has a reduced expression of the form  $stu$  or  $uts$ , where  $m(s, t) \geq 3$  and  $u \in W$ .

## Theorem (CEGKM)

*In type A and B,  $w \in W$  is T-avoiding iff  $w$  is a product of commuting generators.*